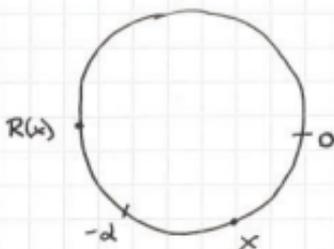


# ① Abelian Powers and Repetitions in Sturmian Words

## Sturmian Words

- Codings of irrational rotations



- $d \in [0, 1]$  irrational
- $R: x \mapsto \{x+d\}$
- $I_0 = [0, 1-d)$ ,  $I_1 = [1-d, 1)$
- $\mathcal{L}(d) =$  language of St words of slope  $d$
- $x =$  intercept

$$\triangleright x_d = 10\dots$$

- Each  $n$ -letter factor  $w = a_0 \dots a_{n-1}$  corresponds to an interval  $[w] = I_{a_0} \cap R^{-1}(I_{a_1}) \cap \dots \cap R^{-(n-1)}(I_{a_{n-1}})$
- $\triangleright x_d$  begins with  $w$  iff  $x \in [w]$

## Continued Fractions

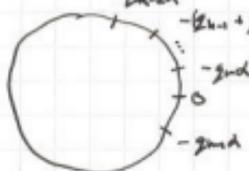
$$\bullet d = [0; a_1, a_2, \dots] = \frac{1}{a_1 + \frac{1}{a_2 + \dots}}, \quad a_i \in \mathbb{Z}_+$$

- convergents  $p_k/q_k = [0; a_1, \dots, a_k]$ ; they satisfy

$$\begin{aligned} p_0 &= 0, & p_1 &= a_1, & p_k &= a_k p_{k-1} + p_{k-2} \\ q_0 &= 1, & q_1 &= a_1, & q_k &= a_k q_{k-1} + q_{k-2} \end{aligned}$$

- semiconvergents  $\frac{p_{k,l}}{q_{k,l}} = [0; a_1, \dots, a_{k-1}, l] = \frac{l p_{k-1} + p_{k-2}}{l q_{k-1} + q_{k-2}}$  with  $1 \leq l < a_k$

- convergents give "best possible rational approximations"; semiconvergents "second best" <sup>to  $d$</sup>



Best in the sense that  $\|q_k d\|$  is smallest, where  $\|x\| = \min\{\{x\}, 1 - \{x\}\}$ .

- $Q_d^- =$  denominators of convergents of  $d$
- $Q_d^+ =$  conv. & semiconvergents  $-a$

## ② Powers and Periods

### Thm (Damanik-Lenz)

- $w^2 \in \mathcal{L}(\alpha)$ ,  $w$  primitive  $\Rightarrow |w| \in Q_\alpha^+$ ,
- $|w| \in Q_\alpha^+ \setminus Q_\alpha$   $\Rightarrow \exp(w) \leq 2$ ,
- $|w| = p_k$   $\Rightarrow \exp(w) \leq a_{k+1} + 2$
- supremum of exponents  $< \infty \Leftrightarrow \alpha$  has bounded partial quotients

0100100 period 3, "covered" by power  $(010)^3$

Thm (Curie-Saari) Minimal period of  $w \in \mathcal{L}(\alpha)$  in  $Q_\alpha^+$ .

### Abelian Powers and Repetitions

- Goal: generalise above theorems to abelian setting.
- Joint work with Fici, Langier, Lecroq, Lefebvre, Mignari, Prieur-Gaston.
- TCS 2015/6 with the same title.
- abelian powers = powers where permutation of letters is allowed

010 · 100 · 001, abelian power of period 3

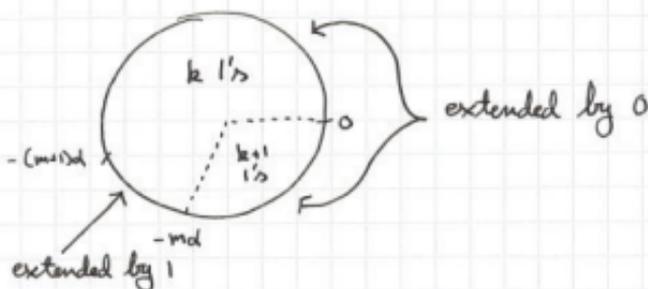
- two words are abelian equivalent if they are permutations of each other
- $w$  is an abelian repetition of period  $m$  if it can be covered by an abelian power of period  $m \rightarrow$  generalises period to abelian setting

01 · 010 · 100 · 001 · 1, abelian repetition of period 3

- ③
- $w$  has abelian period  $m$  if it is an abelian repetition with period  $m$
  - abelian powers harder to study in general, but relatively easy in St words thanks to the following

Thm (Fici et al.) The point  $\{-md\}$  separates the intervals of factors of length  $m$  into two abelian equivalence classes. "heavy" factors above  $\{-md\}$ , "light" below.

Proof sketch. By induction on  $m$ . Case  $m=1$  clear. We prove the case  $\{-(m+1)d\} < \{-md\}$ ; the other case is similar.



Overall:

- below  $\{-(m+1)d\}$ :  $k$  '1's
- between  $\{-(m+1)d\}$ ,  $\{-md\}$ :  $k+1$  '1's
- above  $\{-md\}$ :  $k+1$  '1's

□

- Thus, if  $\{x\}$  and  $\{x+md\}$  are on the same side of  $\{-md\}$ ,  $\{x, d\}$  begins with an abelian ~~power~~ square of period  $m$  etc.
- In general: We have to see how many times  $\|u\|$  fits into  $\| [0, -md) \|$  and  $\| [-md, 1) \|$  to find the max. exponent  $de(m)$  of period  $m$ .

Thm (F. et al.)  $de(m) = \lfloor \frac{1}{\|u\|} \rfloor$ .

- (4) • If  $m \in \mathbb{Q}_d^+$ , then  $\mathcal{L}_e(m)$  is large. Indeed, it can be computed that  $\mathcal{L}_e(g_n) \geq g_{n+1}$ .

$$f = \overbrace{01001010010010100101} \dots$$

- We can make  $\|m\|$  arbitrarily small, so

Thm (Richomme-Scari-Zamboni) Sturmian words contain abelian powers of arbitrarily high exponent.

- This follows from the more general theorem

Thm (R.-S.-Z.) If an infinite word has bounded abelian complexity, then it contains abelian powers of arbitrarily high exponent.

- Since the exponent is always arbitrarily high, it does not make sense to study the max exponent.
- Instead we measure the max ratio between the exponent and the period of an abelian power.

abelian critical exponent:  $\mathcal{L}_c(\alpha) = \limsup_{m \rightarrow \infty} \frac{\mathcal{L}_e(m)}{m}$   
 $= \limsup_{k \rightarrow \infty} \frac{\mathcal{L}_e(g_k)}{g_k}$

Thm (Fici et al.)  $\mathcal{L}_c(\alpha) = \limsup_{k \rightarrow \infty} ([a_{n_1}, a_{n_2}, \dots] + [0, a_n, a_{n-1}, \dots, a_1])$   
 when  $\alpha = [0; a_1, a_2, \dots]$ . Moreover, the following are equivalent:

- $\mathcal{L}_c(\alpha) < \infty$
- $\mathcal{L}(\alpha)$  has bounded powers,
- $\alpha$  has bounded partial quotients.

- Observation:  $\mathcal{L}_c(\alpha) =$  Lagrange constant of  $\alpha$ .
- $\mathcal{L}$  constant of  $\alpha =$  largest  $c$  s.t.  $|\alpha - \frac{p}{q}| < \frac{1}{cq^2}$  for infinitely many  $\frac{p}{q}$

(5)

Abelian Periods

Thm (F. et al) The min. abelian period of any factor of the Fibonacci word is a Fibonacci number.

- Proof idea: Say  $F_k < m < F_{k+1}$  ( $F_k = k^{\text{th}}$  Fibonacci number) and that  $m$  is an abelian period of  $w \in \mathcal{L}(d)$ . (not necessarily min.). Then  $|w| \leq (\beta(m) + 2)m - 2$ . Near the beginning of  $w$  we can always find an abelian power of period  $F_k$  and exponent  $F_{k+1} + F_{k-1} - 3$ . This is sufficient to cover  $w$ .  $\square$

- The abelian periods are typically small:

Thm (F. et al) Let  $k \geq 3$ . The min. abelian period of the finite Fibonacci word  $f_k$  is the Fibonacci number  $n$ , where

$$n = \begin{cases} \lfloor k/2 \rfloor, & \text{if } k \equiv 0, 1, 2 \pmod{4} \\ \lfloor k/2 \rfloor + 1, & \text{if } k \equiv 3 \pmod{4}. \end{cases}$$

- Generalisation "min. abelian period of a word in  $\mathcal{L}(d)$  is in  $\mathcal{Q}_d^+$ " does not work.
- Counterexample in all other cases than Fibonacci.
- Correct formulation:

Thm (P.) Let  $m$  be min. abelian period of  $w \in \mathcal{L}(d)$ . Then  $m \in \mathcal{Q}_d^+$  or  $m = t g_k$  with  $1 \leq t \leq a_{k+1}$ .

- Proof idea: Say  $k$  max. s.t.  $g_k < m$ . The word  $w$  must contain the singular factor of length in all phases mod  $g_k$ . Thus  $|w| \geq (g_k - 1)g_{k+1} + g_k$ ,  $g_{k+1}$  min. return time of singular factor. On the other hand  $|w| \leq (\beta(m) + 2)m - 2$ . These inequalities hold only if  $\beta(m)$  is large i.e.  $d$  is small. Thus  $m$  corresponds to a good rational approximation of  $d$ .